Writing a fast OpenRISC emulator in JavaScript – fun and pain

Sebastian Macke

Vancouver

bears

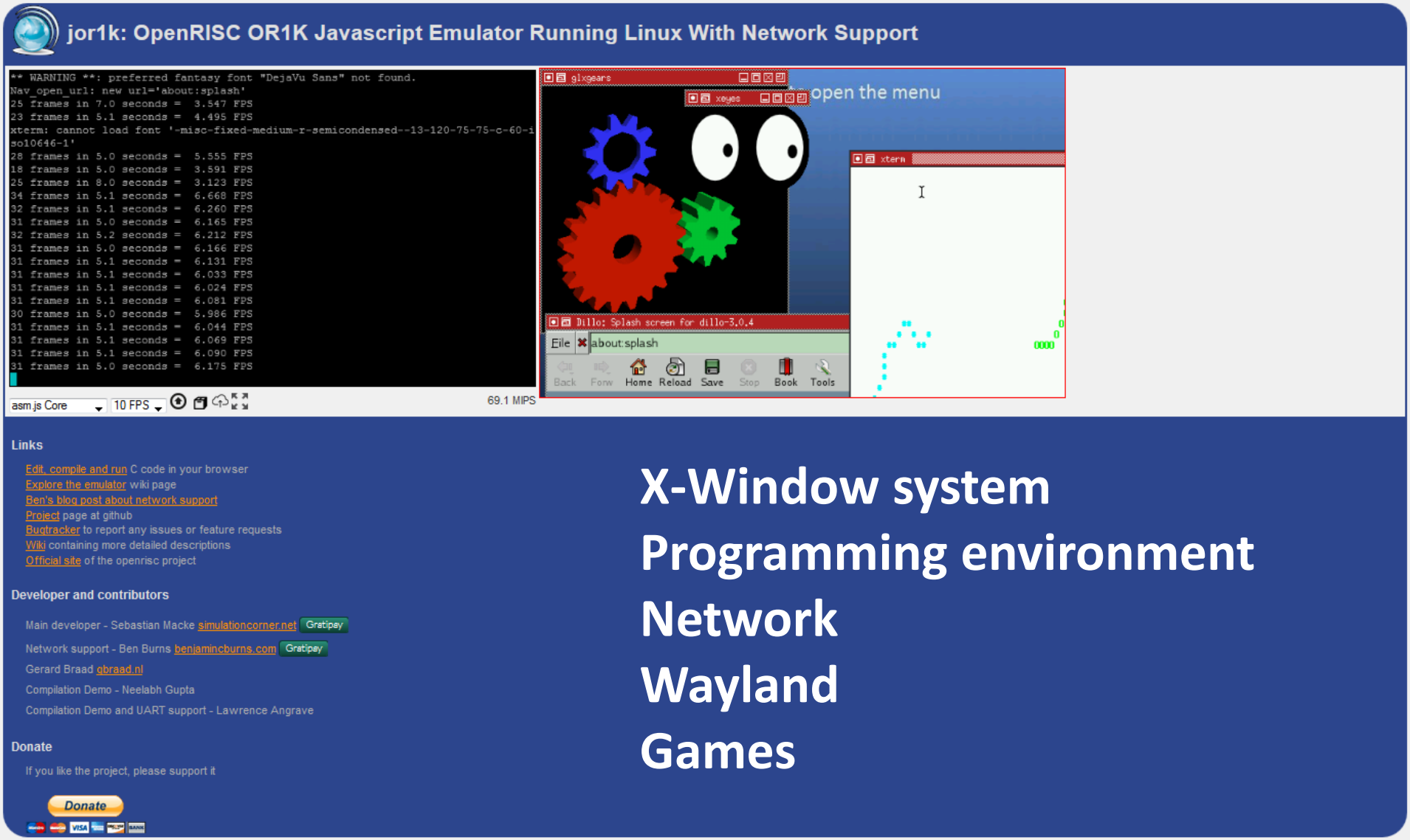
Downtown

No bears

University of British Columbia

Wreck beach
Very liberal

The website: **jor1k.com**



X-Window system

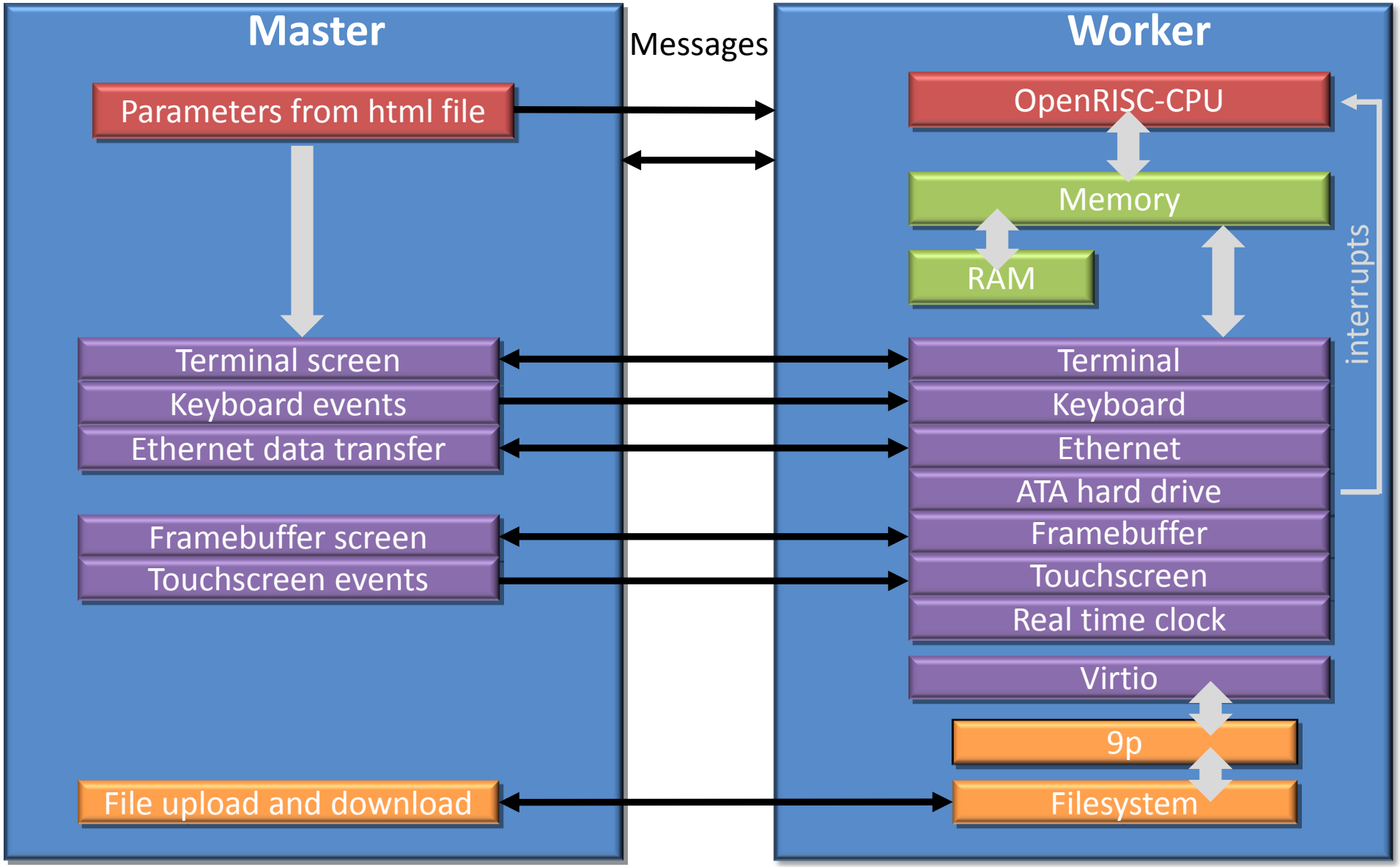
Programming environment

Network

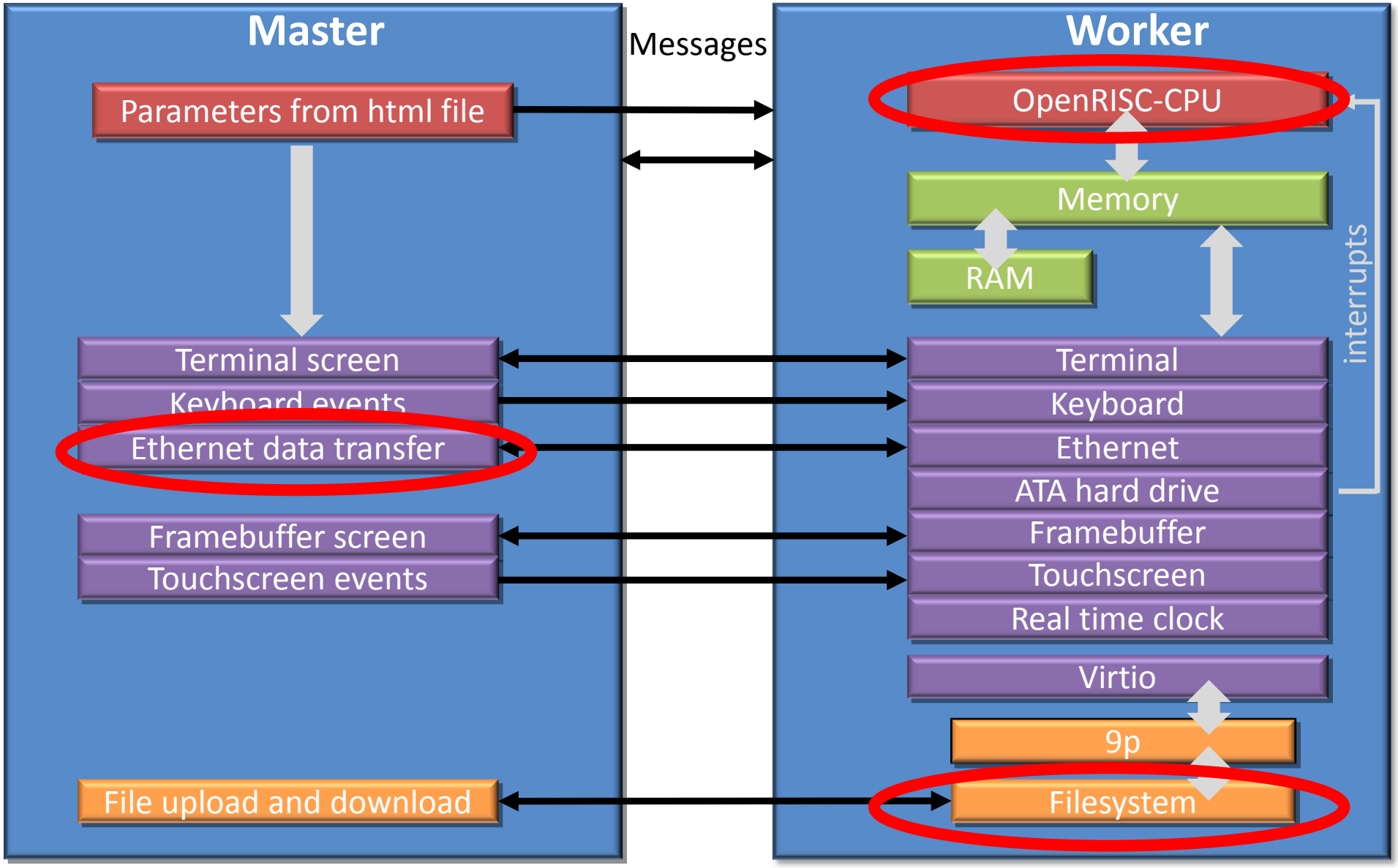
Wayland

Games

JavaScript Modules



JavaScript Modules



Why JavaScript?

- **It runs everywhere**
- Everything which takes more than one click to show does not get much attention.
- It is considered as a bad language, which is sadly true.
  - `[1,2,3]+[4,5,6] => 1,2,34,5,6`
  - `0 == "" => true`
- But the language is better than its reputation.
- At least four companies are writing optimized compilers to squeeze out the maximum performance.

Javascript is not typed

- There are no integers, only doubles, but the compilers try to optimize it
  - `y = 9999999999999999 => y = 10000000000000000`
    (double) (double)
  - `y = 0xFFFFFFFF+1 => y = 0x100000000`
    (integer) (deoptimized into double)
  - `y = 0x7FFFFFFF+1 => y = 0x80000000`
    (Integer) (Int? double?)

- But there are logical operations
  - `y = 0xFFFFFFFF|0 => y = -1`
  - `y = 0xFFFFFFFF>>> 0 => y = 0xFFFFFFFF`

- But there are typed arrays
  - `x = new Uint32Array(length)`

- Ways to optimize:
  - Write like it would be a typed language
  - Take care, that deoptimizations to doubles never occur
    - `y = (0xFFFFFFFF+1)|0 => y= 0x0`
      (integer) (integer)

What is asm.js

- The mode `"use strict";` adds restrictions to JavaScript like additional error messages for accessing undefined variables.
- The mode `"use asm";` adds additional error messages to give you a guarantee for typed variables that must be compiled only once.
  - Only a subset of Javascript is allowed
  - Fully compatible
- Implemented in Firefox in 2013

What is asm.js

Why just error messages?

Firefox with asm:	75.5 MIPS
Firefox without asm:	58.1 MIPS
Chrome (no support for asm):	60.7 MIPS
IE 11 (no support for asm):	68.3 MIPS
Safari on iPAD:	81.0 MIPS

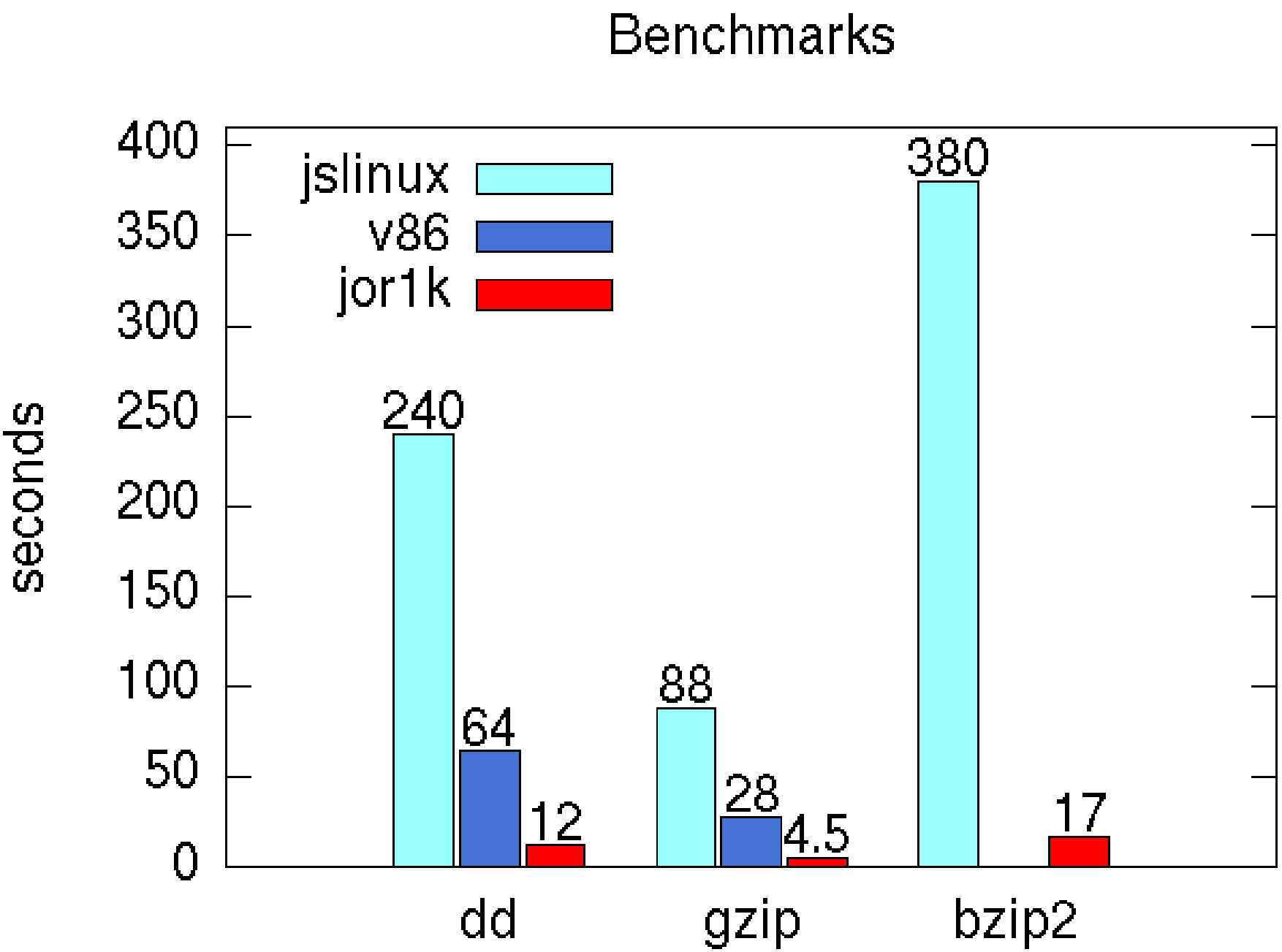
- Implemented in Firefox in 2013

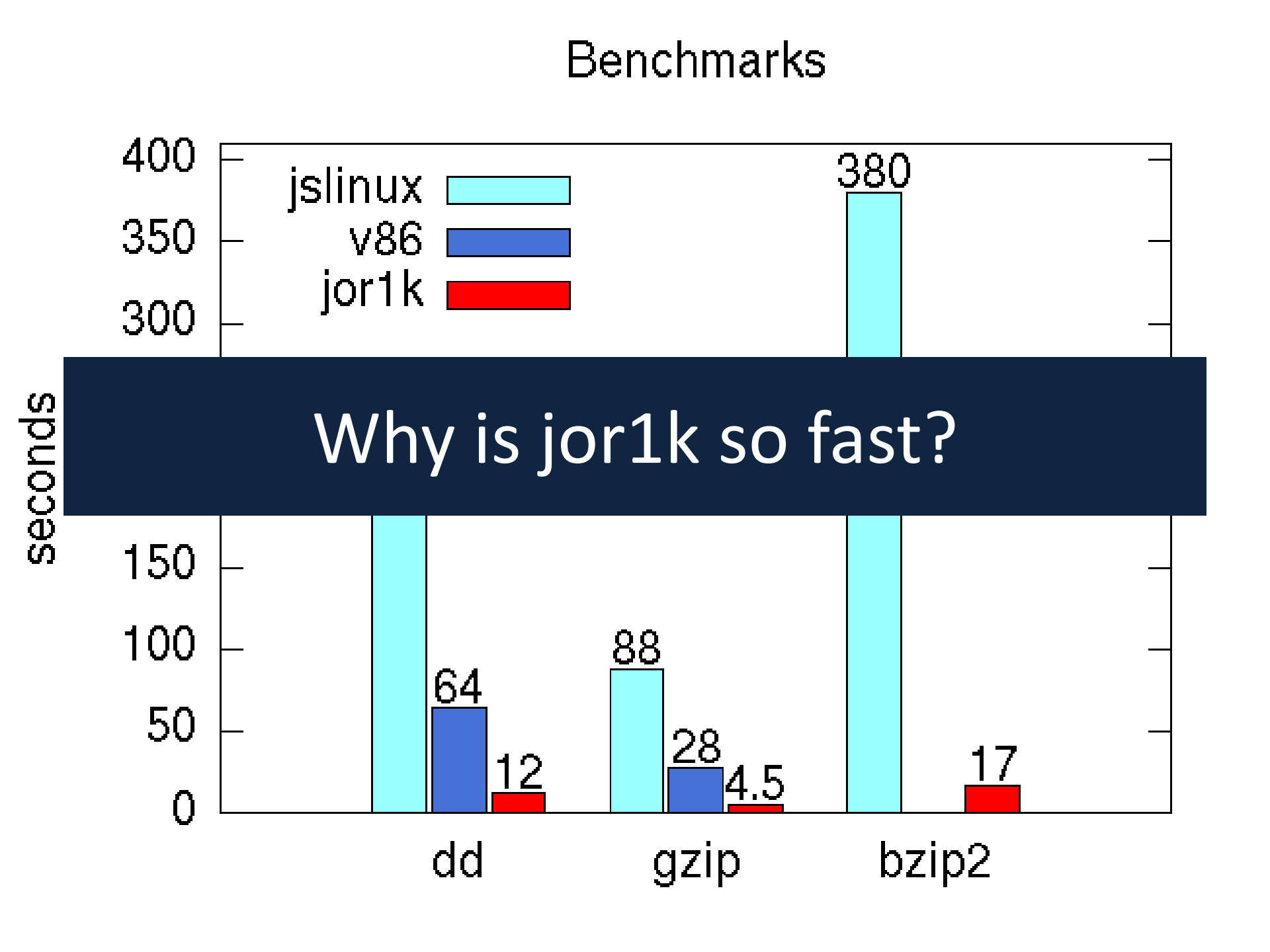
What is asm.js

- But the syntax is nasty
  - `group0[SPR_IMMUCFGR] = 0x18;`

  - `h[group0p + (SPR_IMMUCFGR<<2) >> 2] = 0x18|0;`
- `h` is the heap and `group0p` is the pointer to the table
- In this case the “view” of the heap is 32 Bit. Therefore the last operation for the index must be “>> 2”
- Project Emscripten allows to translate C++ to asm.js JavaScript
  - Switch-case is used instead of goto

The CPU


Benchmarks
seconds
400
350
300
250
200
150
100
50
0
jslinux
v86
jor1k
240
64
12
88
28
4.5
380
17
dd
gzip
bzip2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Why is jor1k so fast?

Architecture

- OpenRISC is easy
  - No history
  - Almost no side effects

```
switch ((ins >> 26)&0x3F) {
…
    case 0x29: // l.andi
        r[(ins >> 21) & 0x1F] = r[(ins >> 16) & 0x1F] & (ins &
0xFFFF);
    break;
…
}
```

=> The instruction set is more like bytecode.

Instruction emulation for ARM

```
void
armv5_and(){
                uint32_t icode = ICODE;
                int rn,rd;
                uint32_t cpsr=REG_CPSR;
                uint32_t Rn,op2,result;
                uint32_t S;
                if(!check_condition(icode)) {
                                return;
                }
                rd=(icode>>12)&0xf;
                rn=(icode>>16)&0xf;
                Rn=ARM9_ReadReg(rn);
                cpsr&= ~(FLAG_N | FLAG_Z | FLAG_C);
                cpsr |= get_data_processing_operand(icode);
                op2 = AM_SCRATCH1;
                result=Rn&op2;
                ARM9_WriteReg(result,rd);
                S=testbit(20,icode);
                if(S) {
                                if(!result) {
                                                cpsr|=FLAG_Z;
                                }
                                if(ISNEG(result)) {
                                                cpsr|= FLAG_N;
                                }
                                if(rd==15) {
                                                if(MODE_HAS_SPSR) {
                                                                SET_REG_CPSR(REG_SPSR);
                                                } else {
                                                                fprintf(stderr,"Mode has no spsr in line %d\n",__LINE__);
                                                }
                                } else {
                                                REG_CPSR=cpsr;
•                               }
•               }
•               dbgprintf("AND result op1 %08x,op2 %08x, result %08x\n",Rn,op2,result);
•       }
```

Neglecting unused features

- CPU flags are not used
- Unaligned memory accesses are not checked
- Snoop hit never happens

> The reservation for a subsequent **l.swa** is cancelled if another store to the same memory location occur, another master writes the same memory location (snoop hit), another **l.swa** (to any memory location) is executed, another l.lwa is executed or a context switch (exception) occur.

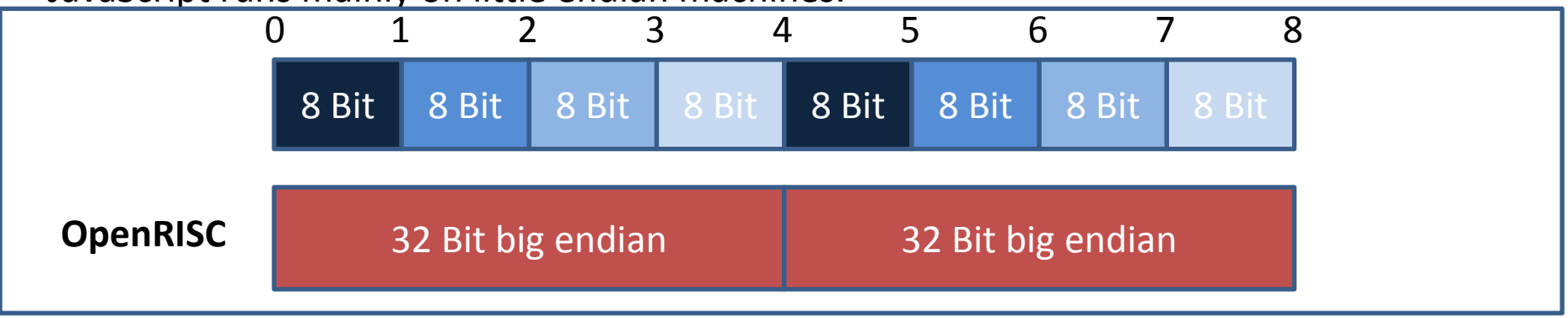
This would add an additional check to the load and store instructions.

Are there any downsides of the architecture to write a performant emulator?

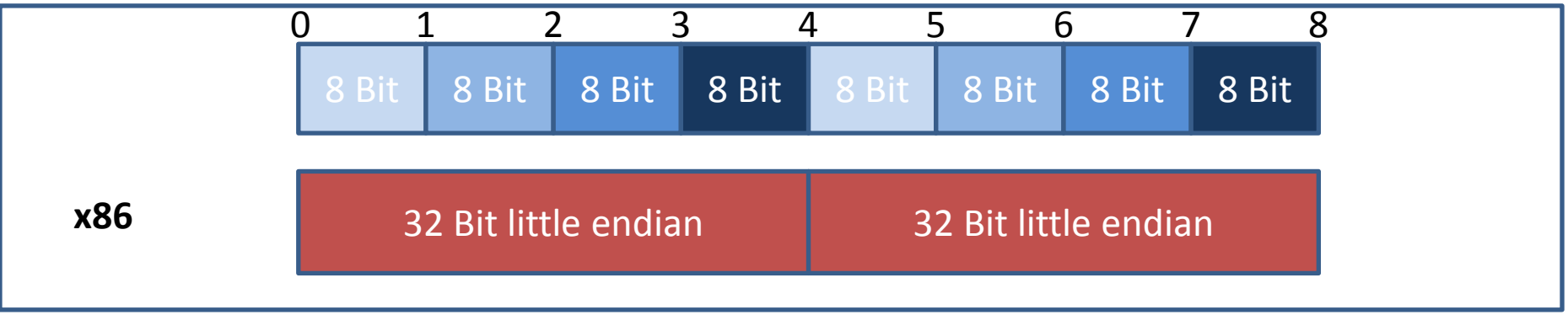
BIG endian on little endian machines

JavaScript allows different views of typed arrays:

JavaScript runs mainly on little endian machines.



But all memory accesses in the emulator are aligned and 32 Bit. So flip every 32-Bit word.



Correction table For little endian machines

Memory access bits	Access
32	`w[addr]`
16	`h[addr^2]`
8	`b[addr^3]`

w[]: 32-Bit view of RAM

h[]: 16-Bit view of RAM

b[]: 8-Bit view of RAM

The MMU: Software TLB lockup

- Usually implemented in two stages
  1. Full translation table in memory
  2. Small translation lookaside buffer (TLB) in the CPU
     - Usually filled in software => But code part translated to JavaScript
- Add third stage
  3. TLB variables (tlb buffer with one entry)

Translation fastpath of virtual to physical addresses:

```
if ((tlb_check ^ virtual_addr) >> 13)
{
    ...
    tlb_check = ...
    tlb_trans = ...
}
physical_addr = tlb_trans ^ virtual_addr;
```

Overhead of the delayed instruction

- Usual instruction pointer increment command line: `"pc += 4;"`
- With delayed instruction support this would turn into
  - ```
    pc = nextpc;
    nextpc += 4;
    ```

- But currently the fastpath for one instruction looks like this:

  - ```
    for(;;) {
      if (ppc == fence) {
        ....
      }
      ins = int32ram[ppc >> 2];
      ppc = ppc + 4;

      switch ((ins >> 26)&0x3F) {
        ....
      }
    }
    ```
  - The idea here is that the virtual pc is computed only when needed by translating `ppc` (physical pc) back to the virtual pc address. The variable `fence` is used to break out of the fast path when `ppc` reaches a jump or the end of the current page.
  - The delayed instruction gives not additional overhead

The Filesystem

How to implement an efficient filesystem with a size of 200MB and 5000 files that runs over the internet?

The Filesystem

- How long does a "`du /`" take over the internet?
  - NFS
  - Samba
  - Sshfs
  - On demand block device

Grab a cup of coffee

Problem is mainly latency, not throughput

Advantages of our filesystem:

- Read only filesystem on server
- Only one user

The Filesystem

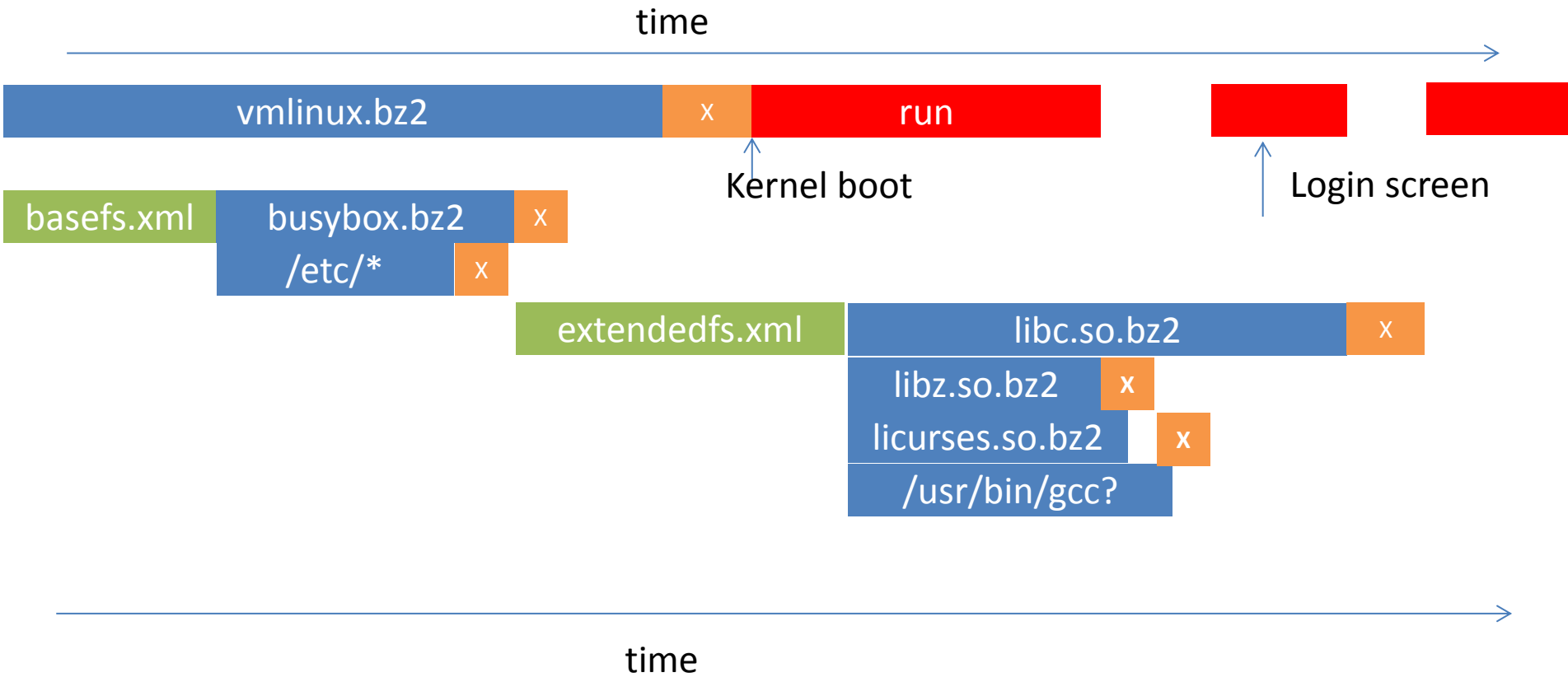
- Implement filesystem outside of the emulator
  - tmpfs like. Use virtio/9p to exchange commands with Linux
- Load the filesystem layout and metadata during the Linux boot process.

```
<Dir name='hp_vndr' mode='40755'>
        <File name='us' mode='100644' size='3339'/>
</Dir>
<File name='ee' mode='100644' size='4167'/>
<File name='ara' mode='100644' size='13194' compressed='1'/>
<File name='ua' mode='100644' size='14943' compressed='1'/>
```

- Load compressed files on demand.
  - OpenRISC binaries compress really well
  - .bz2 currently, in future .xz
  - Ordinary web server needed

- Future: dependencies between files, packages
  - http 2.0 will help here

Booting process timeline



Additional features of the filesystem

- Upload files into home folder
- Download home folder (as .tar)
- Sync with server
  - Unique user id (`http://s-macke.github.io/jor1k/?`**`user=cdqKKPxjfa`**)
  - Currently 1MB quota
  - server only needs upload.php

Network

- *Yo dawg, I heard you like browsing the web, so I put a browser in your browser so you can browse while you browse!* (Twitter user Scott Elcomb)

Network

- Server in the USA
  - connected via websockets
  - Sending and receiving ethernet frames connected to a Linux TAP device
- Full working intranet
  - Start jor1k in two windows and open a ssh session between them.
- Major network applications available
  - wget, curl, nc, ping, traceroute, telnet, ssh, nmap
  - Openssl with certificates
  - Web browsers: lynx, links, dillo

Future

- Sound (implemented but not activated)
- SMP
- Run Debian (just one bug left)
- Run Firefox (70% compiles)
- Status, statistics and debug screen
- Download already booted Linux (state file)
- More terminals, better user interface, direct access to the filesystem tree.
- Dynamic recompiler with the `eval` function?

Suggestions welcome

Thanks

- **Stefan Kristiansson** for the toolchain and infinite help in the chat.
- **Ben Burns** for implementing the network and providing the relay server
- **Lawrence Angrave and Neelabh Gupta** for the C-development website
- **Jonas Bonn** for the Linux kernel support
- **Christian Svensson** for the OpenRISC Debian distribution


Play Monkey Island
MONKEY ISLAND


Develop in C
// Compile this program to find the syntax
#include <stdio.h>
int main() {
printf("Hello World!\n")
return 0;
}
Compile and run


Play Elite II
Mouse
Actual:0.0 kmhr⁻¹
Set:0.0 kmhr⁻¹
Alt.0 m
Relative to:Merlin
12:00:03 1-Jan-3200
Landed
External View

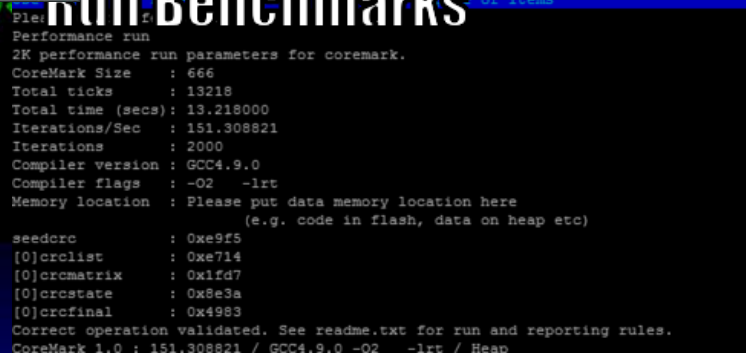

Play DOOM
47
AMMO
100%
HEALTH
ARMS


X Window system available
glxgears
to open the menu
xeyes


Browse
Dillo: Google
http://www.google.com/
Back
Save
Stop
Book
Tools
Search Images Maps Play YouTube News Gmail Drive More »
Web History | Settings | Sign in
Google
Advanced search


Watch movies


Play Toppler


Run Benchmarks
Performance run
2K performance run parameters for coremark.
CoreMark Size : 666
Total ticks : 13218
Total time (secs): 13.218000
Iterations/Sec : 151.308821
Iterations : 2000
Compiler version : GCC4.9.0
Compiler flags : -O2 -lrt
Memory location : Please put data memory location here
(e.g. code in flash, data on heap etc)
seedcrc : 0xe9f5
[0]crclist : 0xe714
[0]crcmatrix : 0x1fd7
[0]crcstate : 0x8e3a
[0]crcfinal : 0x4983
Correct operation validated. See readme.txt for run and reporting rules.
CoreMark 1.0 : 151.308821 / GCC4.9.0 -O2 -lrt / Heap